Programming the EDSAC



Andrew Herbert The National Museum of Computing 16th November 2022

EDSAC FIRSTS

• The world's first PRACTICAL electronic digital stored program computer = computer of the modern kind

• The world's first computer programming system





MAURICE V. WILKES, F.R.S. DAVID J. WHEELEB STANLEY GILL

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THE PREPARATION OF The Preparation of Programs AN ELECTRONIC DIGITAL COMPUTE for an Electronic Digital Computer

Maurice .V. Wilkes, David .J. Wheeler and Stanley Gill Addison Wesley, 1951

With special reference to the use of the EDSAC

Why a Programming System?

The methods of preparing programs for the EDSAC were developed with a view to reducing to a minimum the amount of labour required, and hence of making it feasible to use the machine for problems which require only a few hours of computing time as well as for those which require many hours. This necessitated the establishment of a library of subroutines and the development of systematic methods for constructing programs with their aid.

[WWG 1951]

Note emphasis on programmer productivity rather than on "optimal programming".

To the potential user of an automatic digital calculating machine, the successful design and construction of the machine itself is only a first step, though certainly an essential one. In order that the machine should in practice be useful to him in the calculations he may desire to carry out with its aid, the provision of an adequate organization for using the machine is as important as the machine itself.

The process of building up [such] a library of subroutines, and testing its value by practical use, appears to have proceeded further at the Mathematical Laboratory of the University of Cambridge than elsewhere.

... it is a practical and useful system has been tested by experience; it divests programming of the appearance of being something of a magic art, closed except to a few specialists, and makes it an activity simple enough to be undertaken by the potential user who has not the opportunity to give his whole time to the subject.

The subject is one which is still developing. . .

- 1. Store: ultrasonic delay line holding 1024 ★ 17 binary digit numbers stored in true two's complement form, so most significant digit corresponds to the sign of the number.
- 2. Numbers are held in 1024 numbered "storage locations" numbered serially from 0 to 1023 for reference. Each such reference number is often called the "address" of the associated storage location.
- 3. Hence 17 bit numbers are often called "short numbers"
- 4. Two consecutive storage locations, starting from an even numbered address can be combined to make a 35 binary digit "long number".
- 5. Arithmetic unit: add, complement, collate, shift. Combine to enable subtract, multiply, round, but no divide.
- 6. Accumulator register of 71 bits.
- 7. Multiplier register of 35 bits.
- 8. Input: 5 hole paper tape read by photoelectric reader.
- 9. Output: teleprinter.
- 10. Control: an "order" passes from store into the control unit ("Stage I"), then it is executed {Stage II"). The machine then, generally, then automatically takes the next order from the location following that of the order just executed.

EDSAC Block Diagram



Order Code: O A F or O A D O: Function code – operation to be carried out A: Address of location to be used as operand F or D: Short or Long number

- A n Add C(n) to Acc
- S n Subtract C(n) from Acc
- H n Copy C(n) to Multiplier
- V n Multiply C(n) by C(Mult) and add product to Acc
- N n Multiply C(n) by C(Mult) and subtract product from Acc
- T n Transfer C(Acc) to location n and clear Acc
- U n Transfer C(Acc) to location n but do not clear Acc
- C n Collate C(n) with C(Mult) and add to Acc
- R D Right shift Acc one place ($\times 2^{-1}$)
- L D left shift Acc one place (× 2)

L 2 ^{p-2} F	Multiply by 2^{p} (2 < p <= 12)
R 2 ^{p-2} F	Multiply by 2 ^{-p} (2 < p <= 12)
EnF	If C(Acc) >= 0 execute next order from location n; otherwise proceed serially
G n F	If C(Acc) < 0 execute next order from location n; otherwise proceed serially
l n F	Read next 5 bit code from input to location n from tape reader
O n F	Print character set up on teleprinter, then set up m.s. 5 bits of location n as next character
F	Read back last set character
Х	Ineffective – no-op
Y	Round Acc to 34 digits (i.e., add 2 ⁻³⁵).

Z Halt and ring the bell

Fixed Point Arithmetic

Binary point assumed between top two most significant bits, so numbers are -1 <= x < 1.

Thus A order computes x+y-2 if x+y>=1, and x+y+2 if x+y<-1.

When two long numbers are multiplied together the resulting 69 digits are available in the Accumulator.

0.5 + 0.25 + (0.5 * 0.25) = 0.875

Compute x+y+xy; x in location 6, y in location 7.

- (0) T 8 F
- (1) A6F
- (2) A 7 F
- (3) H 6 F
- (4) V7F
- (5) Z O F
- (6) +0.5 0100000000000000
- (7) +0.25 0010000000000000
- (8) (spare)

Demo1

Integer Arithmetic

Can treat accumulator as holding integers for A, S, C, L, R, E, G but for N, V multiplier is always treated as a fraction.

i.e., integers are stored as value * 2⁻¹⁶ so need to multiply by 2¹⁶ after multiplication.

10 + 5 + (10 * 5) = 65

Compute x+y+xy, x in location 8, y in location 9.

(0)	T 10 F	
(1)	H 8 F	
(2)	V 9 F	
(3)	L 512 F	; multiply by $2^{16} = 2^{11*}2^5$
(4)	L 8 F	
(5)	A 8 F	; note addition after
(6)	A 9 F	; multiplication
(7)	Z 0 F	
(8)	+10	000000000001010
(9)	+5	000000000000101
(10)	(spare)	

Demo2

Loops

- Loop to print digit 7 five times
- C(11) is "figure shift"
- C(12) is "7"
- C(13) is RET
- Loops while Acc < 0 (-5, -4, -3, -2, -1)
- Note: need to set figure vs. letter shift
- Note: output delayed one character
- Note: only G < 0 and E <= 0, but no "equals" order

(0)	T14F	
(1)	O10F	; Type figure shift
(2)	S9F	; Set count = -5
(3)	A10F	; Increment count
(4)	012F	; Type '7'
(5)	G3F	; loop if count < 0
(6)	013F	; Type RET
(7)	011F	; Type figure shift
(8)	ZOF	; Halt
(9)	+5	
(10)	+1	
(11)	0101100	0000000000
(12)	0011100	0000000000
(13)	1100000	0000000000
(14)	(spare)	

Demo3

Indexing

•	Initially EDSAC had no index register.	
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- Invented for Manchester Mark 1 in 1949.
- Later adopted by EDSAC.
- So to do an indexed calculation, e.g., sum a vector, we have to write self-modifying code that manipulates program in store.
- To do arithmetic on orders we need to understand binary format of orders:
 - 5 m.s. bits: order code
 - 1 bit : spare (later add B register to address)
 - 10 bits: address in range 0-1023
 - I.s. bit: 0 = F, 1 = D
- Demo adds contents of vector starting at location 15
- We have to add +2 to location 2 each time around the loop to fetch the next element of table

U U		
(0)	T13F	; Clear
(1)	A14F	; Add s
(2)	A15F	; Add y
(3)	T14F	; Store
(4)	A2F	; C(2)
(5)	A12F	; C(2)+
(6)	U2F	; Mod
(7)	S20F	; Chec
(8)	G0F	; Loop
(9)	T13F	; Clear
(10)	A14F	; Resu
(11)	ZOF	; Halt
(12)	+2	; addr
(13)	+0	; work
(14)	+0	; sum
(15-19)	+1 +2 +3 +4 +5	; vecto
(20)	A20F	; senti

' acc sum vector[0] ** e in sum +2 (inc. address) lify (2) ** ck for end if not at end r acc ılt ess stride space or (Demo4) inel

Emulators etc

- Demos thus far run using EDSAC team test program generator:
 - Assembler written with "modern" facilities
 - Emulator written in C with tracing facilities
 - adapted from original by Lee Wittenberg
 - On EDSAC we use a "Signal Sequence Injector" box to set up program in main store from location 0 onwards
- Visit GitHub/andrewjherbert to find these:
 - edsacasm <u>https://github.com/andrewjherbert/edsacasm</u> Python
 - edsac <u>https://github.com/andrewjherbert/EDSAC-Emulator</u> C
- But this is not how EDSAC users wrote code...

Ale in the	EDSAC PROGRAMME SHEET								
Count of	for the state	,	REF	DATE					
Calculation of a	unes pri yez e	fc,							
Calculates /2) 128 vg) 64 vg								
Use with tap	e WSG :								
Order	Notes		Order	Notes					
• P F	-	this fo	V 2047 D	2 starts at sin 20					
$2 P \times F$	x= 4f x 2048	Res 2	R 4095 D P & (FD)	$h = \sin^2 D \times 32768$					
3 T 126 K	F in 126	23	PF	A starting value.					
5 T 294 K	<u>.</u>	5	P F	S clears 2					
6 E 231 F T 231 K		6	P F. T 258 K						
8 A 231 F	State of the second second	8	A 243 F	1					
9 4 165 F		9	T 126 K T 323 K						
1 T 288 D		1	T 171 E						
2 A 255 F 3 G 91 F		3	/E 179 F						
4 A 288 D		4	T 179 K	1					
6 0 241 F		5	T 132 K						
7 E 296 F 8 Ø F	DIF due)	7	P 10813 F P 32000 F						
9 W F -	= (Chang do 1 - 2010 x 32768 to deps)	9	T 317 K						
0 P 542 F 1 T 329 K	$\Delta 1 \sim 2048 = g(x = f).$	0	P 256 F T 314 K	0 241 F T D					
2 A 126 D A 2112 E		2	T 36 D	E 316 F					
4 T 126 D	Sa'= no of steps before	3	E 248 F T 248 K	T 366 K P ~ F					
5 T 355 K	Goest value of sint D	5	H 36 D	T 211 K) to					
7 A 242 F		, 6 7	Y 288 D T D	S DS X					
18 T 128 D		8	A 251 F	E 144 K D F					
1 120 N.	*		7 9/4	F					









 $\begin{array}{c} B = \frac{7}{76} \\ + 032154 + 019084 \\ + 033505 + 018876 \\ + 033505 + 018732 \\ + 035237 + 018687 \\ + 037588 + 018794 \\ + 040916 + 016138 \end{array}$ +040919 +019138 $\begin{array}{r} +040315 + 019138 \\ +045875 + 019864 \\ +053687 + 021222 \\ +066905 + 023656 \\ +091337 + 027966 \\ +141844 + 035461 \\ +268904 + 047536 \end{array}$ 15 +031130 +019063 +032182 +018868 +033528 +018743 +035304 +018723 +037726 +018863 +041163 +019252 +041163 +019252 +054331 +020039 +054331 +021476 +067907 +024009 +092799 +028414 +143687 +035922 +270334 +047789 +031106 +019048 +032179 +018867 +033558 +018867 +033558 +018759 +035380 +018763 +037871 +0189366 +041412 +019366 +046680 +020213 +054959 +021724 +068860 +021724. 10714 ... +068860 +024346 +094165 +028832 +145354 +036339 +271576 +048008 +031051 -035

Initial Orders

Fixed program to load source programs from paper tape into store

Input is alphanumeric

Combined assembler and linker to enable user code to be linked to predefined library routines

Unique to EDSAC

Programming tour de force by David Wheeler



Proc. Royal Society A, 202, August 1950: D.J. Wheeler, Programme organization and initial orders for the EDSAC. https://royalsocietypublishing.org/doi/10.1098/rspa.1950.0121

Initial orders concepts

- Instructions in alphanumeric form rather than binary
 - Like modern assembly code
- Control codes to direct initial orders where to load and how to fix up addresses, start execution
 - To enable linking in subroutines in arbitrary order
- Addressing relative to a previously set parameter (control code)
- But no error handling!

Warwick Simulator

Written by Martin Campbell-Kelly Available for Windows and MacOS GUI replicates original EDSAC operation

🖳 Edsac	
	Output From: OXO
	987 NOUGHTS AND CROSSES
	654 BY 🔤
	3 2 1 A S DOUGLAS, C.1952
	LOADING PLEASE WAIT
	RDSAC/USED RIDST (DIAL 0/1)-0
	DIAL MOUR-7
	DIAL NOUR.C
	DIAL HOVE:8
	Clear Reset Start Stop Single E.P
Constant SCR Constant Order T Multipl	ank ier icand IV Short Tanks Acc

https://edsac.net

Includes Tutorial Guide, original EDSAC subroutine library, worked example programs

Hello world

T 64 K G K and E Z P F are control combinations

 $\boldsymbol{\theta}$ is a "parametric address" T 64 K – load from loc 64 G K – set $\boldsymbol{\theta}$ (to 64) E Z P F – enter program at Location $\boldsymbol{\theta}$ (64)

* Is erase character (32 decimal)

N.B. Data input as instructions

		Т	64 K	Load from location 64	
		G	K	Set θ parameter	
Start \rightarrow	0	Ζ	F	Stop	
ire	1	0	5 θ	Print letter shift	T64K
	2	0	6θ	Print "H"	GK
	3	0	7 θ	Print "I"	ZF
ress"	4	Z	F	Stop	050
54	5	*	F	Letter shift	07@
	6	Н	F	"H"	ΖF
ו at	7	I	F	п ^Т п	*F нг
		E	Z		IF
		Р	न	Enter at location $0 heta$	EZPF

(a) Program text

(b) Program tape

Demonstration Programs/Hello.txt

EDSAC character codes

Programs prepared on perforator Note NO figure / number shift

Output produced in Teleprinter code

Order field in instruction is the bit pattern of the order character, i.e., A = 11101

Note convention for typing Greek letters when using emulators

Table 2 Edsac Character Codes

Perfo	rator	Telep	rinter	Binary	Decimal
Letter shift	Figure shift	Letter shift	Figure shift		
Р	0	Р	0	00000	0
Q	1	Q	1	00001	1
W	2	W	2	00010	2
E	3	Е	3	00011	3
R	4	R	4	00100	4
Т	5	Т	5	00101	5
Y	6	Y	6	00110	6
U	7	U	7	00111	7
Ι	8	Ι	8	01000	8
0	9	0	9	01001	9
J		J		01010	10
р		Figure S	Shift	01011	11
S		S	"	01100	12
Z		Z	+	01101	13
K		Κ	(01110	14
Erase ¹		Letter S	Shift	01111	15
Blank tape	2	(no eff	fect)	10000	16
F		F	\$	10001	17
q		Carriage	Return	10010	18
Ď		D	;	10011	19
f		Spac	e	10100	20
Н	+	н	£	10101	21
N	-	Ν	,	10110	22
М		М	•	10111	23
D		Line F	eed	11000	24
L		L)	11001	25
Х		Х	/	11010	26
G		G	#	11011	27
А		А	-	11100	28
В		В	?	11101	29
С		С	:	11110	30
v		V	=	11111	31

Notes

1 Erase is represented by an asterisk ("*") in the simulator. When this character is *output*, it sets the teleprinter into letter shift.

2 Blank tape is represented by a period ("."). This character has no effect on output.

3 The personal computer text environment has only a "newline" character. On the Edsac simulator, the line-feed character is interpreted as a newline character, and carriage returns are thrown away.

4 The symbols q, f, D or p are typed as @, !, & and #, respectively.

Control combinations

- T m K set load point to m
- G K set θ parameter to load point
- TZ restore **θ** parameter
- E m K P F enter program at location m
- EZPF enter program at location $\boldsymbol{\theta}$
- P Z or P K start of new tape block

Subroutines – the Wheeler jump

A m F to pick up where calling from (m+1) A m F is 11000... so negative

C(3) = U 2 F

Calculate E m+2 F and store as final instruction

Return to caller

m	A	m	F	pick up self	
m+1	G	n	F	jump to subroutine	master routine
m+2		•		control returns here -	
n	A	3	F	form return link	
n+1	Т	р	F	plant return link	
		•			subroutine
		•			
р	()	return link planted here -	

Cubes

Nichomacus' formula for cubes:

$$1^{3} = 1$$

 $2^{3} = 3 + 5$
 $3^{3} = 7 + 9 + 11$
 $4^{3} = 13 + 15 + 17 + 19$
etc

Use library routine P6 to print integers

.. represents blank tape

	31	30	29	28	27	26	25	24	23	22	21	20	19	18	17	16	15	14	21 → 13	12	Ħ	10	9	P6 → 8		6	ы	4	د	22 → 2		Enter→ O	
(a) N	Δ	θ	я	ų	P	Р	Р	P	ъ	R	G 1	A 2	A 2	T 2	A 2	U 2	A 2	A 21	T 2	S 2	T 2	A 2	A 2	T 2	G 5	A	ы	A 2	03	0	0 2	7 G	<u>د</u>
Master	Ē	ч	ㅋ	Ē	D	ы	D	D	Ð	2 0	3 0	7θ	θ θ	3θ	3 (5 0	θ	5 0	θ	4 0	40	7 0	4 0	3θ	6 F	0 0	ы	3 0	1 0	θ	9 9	च >	7
routine	If	CL	figs	=2	Ë	count	m (=1 initially)	n (=1 initially)	k (n ³ ; =1 initially)	Repeat main cycle	Jump to 13 if count • 0		Increment count		7 k+m to k		m+2 to m	,] -n to count		n+1 to n		Zero to k		Print OF using P6]k to OF		New line	Fiqure shift	Ston Det N-haramerer	
																												Master	P6 (print)			Routine	
125 216 512	195 195	27	∞ +			,	(c) Make-up of		EZPF	ſ		Mae+		space r 2		Γ	0.1	۶¢		1 00 1	T 76 K	ahare i i		(0) THOTA (0)	(h) Tahle of			88	56		first order	Location of	
							program tape			L	U F	70			-							_			Pointines				32		locations occupied	Number of storage	
(d) Program tape		1. 	#7		PF	נק	PD	DD TAKE	G13@	A26@ A27@	T230	A230	A290	A25@	7260	224A	H110	A24@	T23@	G56F	160 11	A230	D		ons			LAFT4DALFS3@G9@EFSF031@E20@d995FUFIF	V4DU4DAFG26@TFTF05FA4DF4FS4F	[20] (XA 3PM756H796UTFM1DA 36MFH3065CATAF	ms	 ∕Cu	ıbes.txt

(e) Printout

Conventional "coding sheet" style for writing programs No layout on EDSAC tape No comments on EDSAC tape Use of **θ** to make code position independent Constants written as pseudo orders

Advanced features

Code letters:	Code-letter	Location	Value
	F	41	0
	q	42	Origin of current routine
	Ď	43	1
	f, H, N, M V	44, 45, 46 55	For use by programmer

Used to create position independent code and data cross references

Subroutine parameters:

Pass via fixed address (often 0) Include in calling sequence

Run and delete open subroutines on the fly to save store...

Divide subroutine Divide subroutine State P 7 Print subroutine Main program M block at 21 - data State	 3 θ F carriage return 4 Δ F line feed 5 M F decimal point 6 P F count 7 W F = n (=2·2⁻⁴ initially) 	$\begin{array}{cccccccccccccccccccccccccccccccccccc$	$19 \rightarrow 1 \qquad \text{T} \qquad 6 \text{ M} \qquad 3 \text{ set count to -y}$ $2 \qquad \text{A} \qquad 2 \text{ M} \qquad 1 \cdot 2^{-4} \text{ to } 00$ $3 \qquad \text{T} \qquad 0 \qquad 1 \cdot 2^{-4} \text{ to } 00$ $3 \qquad \text{T} \qquad 0 \qquad 1 \cdot 2^{-4} \text{ to } 0F$ $5 \qquad \text{T} \qquad 4 \qquad \text{A} \qquad 7 \text{ M} \qquad 1 \cdot 2^{-4} \text{ to } 0F$ $5 \qquad \text{T} \qquad 4 \qquad \text{A} \qquad 7 \text{ M} \qquad 1 \cdot 2^{-4} \text{ to } 0F$ $6 \qquad \text{A} \qquad 6 \qquad 6 \qquad 3 \text{ M} \qquad 1 \text{ set } 00 \text{ to } 00/40/\text{ie. } 1/n)$ $6 \qquad \text{A} \qquad 6 \qquad 6 \qquad 3 \text{ Monormal point and point decimal point decimal point decimal point at loc 56}$ $10 \qquad \text{A} \qquad 11 \qquad 6 \qquad \text{Print } 0$ $11 \qquad \text{A} \qquad 11 \qquad 6 \qquad \text{Print } 0$	G K T 47 K Set M parameter P 21 0 T Z
Main program M block at 21 - data 10 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1 1	(c) Make-up of program tape	space P Z P1 Master	Divide subroutine	space PK T 56 K
Index Index <th< th=""><th>(d) Printout</th><th>M block at 21 - data , 1999999999999999999999999999999999999</th><th>Main program Enter main program (0) Table of routing - 21 - 20</th><th>Routine Location of Number of stora first order locations occup</th></th<>	(d) Printout	M block at 21 - data , 1999999999999999999999999999999999999	Main program Enter main program (0) Table of routing - 21 - 20	Routine Location of Number of stora first order locations occup

(a) Master routine

• Debugging – post-mortem

Start reciprocals

... Executes ...

Start PM5

Dial start location, e.g., 134 (113+21 = start of data)

Debugging – Checking (i.e., tracing)

Assemble program with checking routine at end

C7 – execution trace

C10 – arithmetical trace



Using command line emulator

Demo5

punch – convert ASCII to EDSAC code

same conventions as Warwick emulator for special symbols etc

edsac - run emulator taking input from stdin

-v1/-v2 tracing

-Innn order limit

-s to start

-b for EDSAC replica SSI emulation

tprint – convert Teleprinter output to UTF

By contrast...

From Turing's programming guide for Manchester Mark 1



Fig. 2. (continued).

How to get started

- Download Warwick simulator, work through examples
- Pitfalls:
 - Remember the store is tiny
 - Be careful about long versus short numbers.
 - Remember to scale calculations.
 - Remember no index registers so vectors, arrays and stacks tedious to manipulate consider writing subroutines / interpreters
 - Read library subroutine specifications carefully to understand parameter passing conventions and any special control combinations to load them.
 - Use code letters to divide code and data into short blocks to avoid having to renumber addresses if additional code or data inserted (or deleted).
 - Beware miscoding pseudo orders (i.e., constants)
 - Must use library routines (R series) to input long numbers n.b., R2 will input long integers at load time
 - Remember need to set teleprinter shift and to force out last character